



ACM MobiHoc 2020

Optimizing Information Freshness using Low-Power Status Updates via Sleep-Wake Scheduling

Ahmed M. Bedewy[‡]

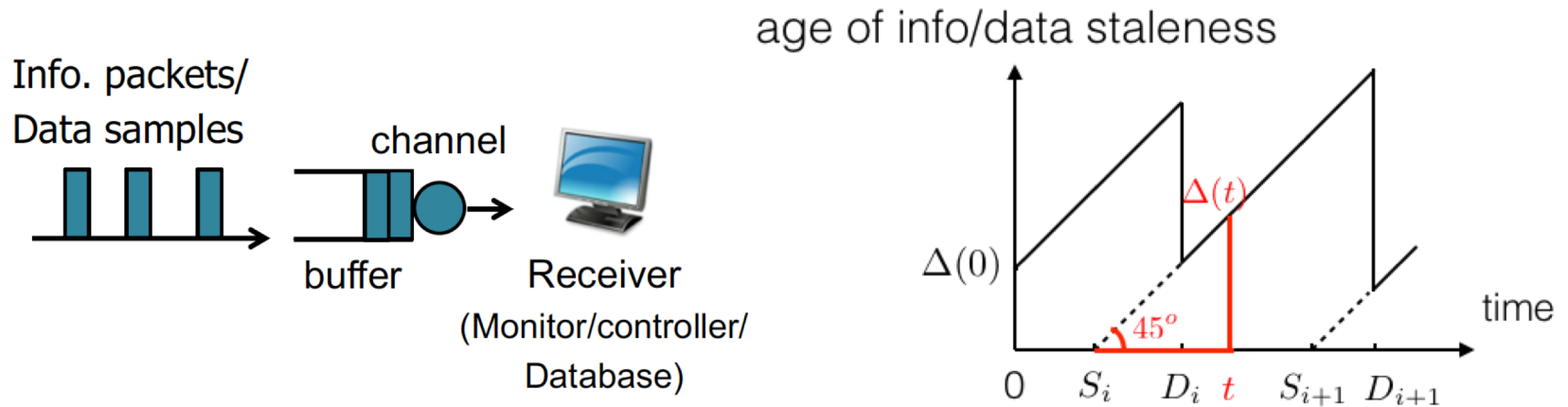
Joint work with Yin Sun^{*}, Rahul Singh[‡], and
Ness B. Shroff[‡]

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Oct. 12th 2020



What is the Age of Information?



Definition: at any time t , the **age-of-information (AoI)** $\Delta(t)$ is the “age” of the **freshest** sample available at the **destination** before time t

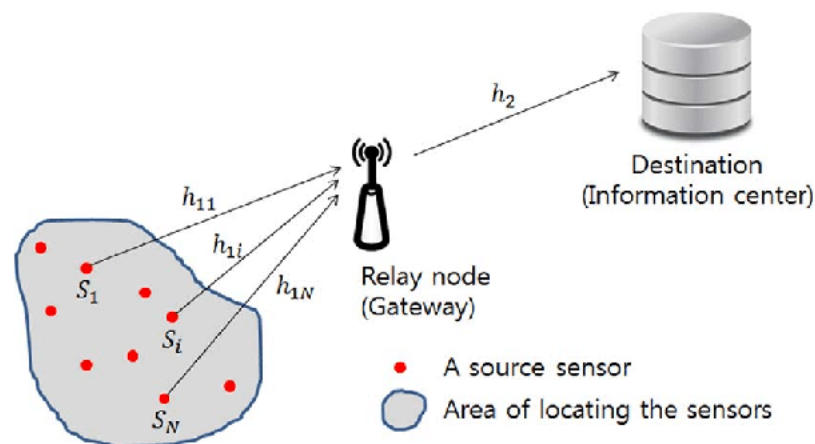
- If sample i is generated at S_i and delivered at D_i

$$\Delta(t) = t - \max\{S_i : D_i \leq t\}$$

- Age **grows linearly**, and **drops** upon new sample delivered

Motivation

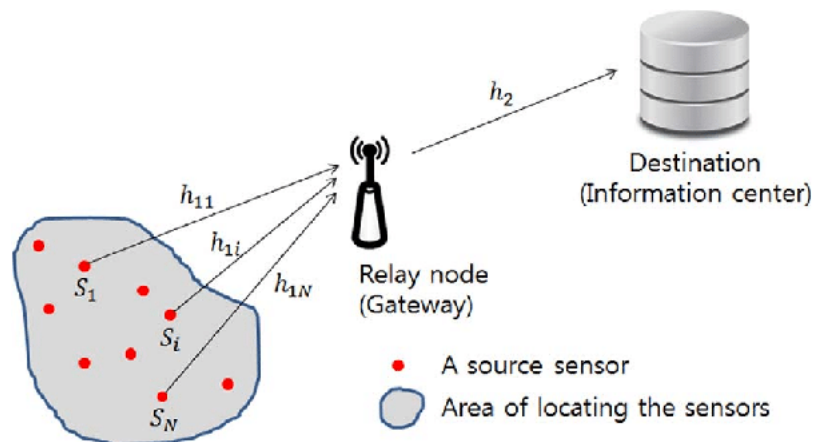
- **Wireless sensor networks**
 - Sensor nodes in **remote** or **hard-to-reach** areas
 - Sharing **same** channel
 - Required to operate unattended for **long durations**.
 - Have **limited** battery capacity



WSN to observe environmental phenomena

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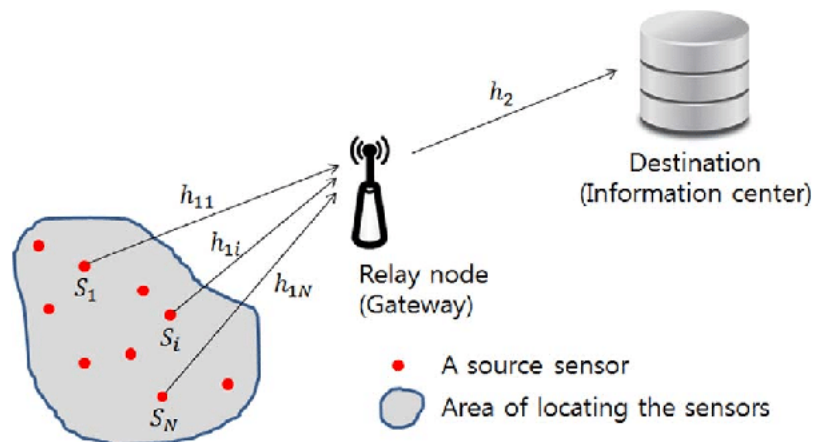
WSN to observe environmental phenomena



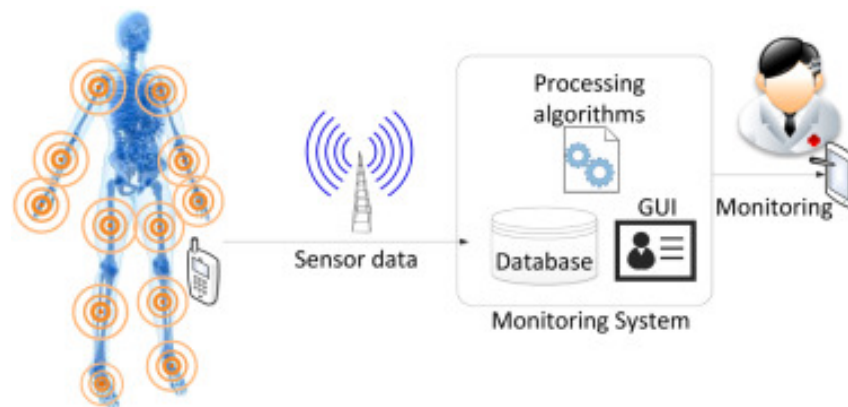
Medical sensor networks

Motivation

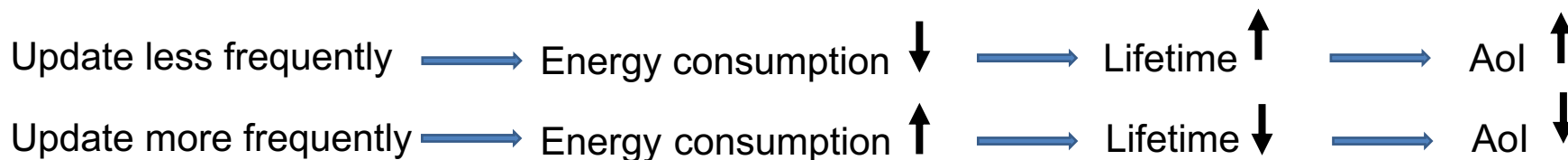
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 - Sensor nodes in **remote** or **hard-to-reach** areas
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 - Maintain **low** Aol
 - **Long** Lifetime



WSN to observe environmental phenomena

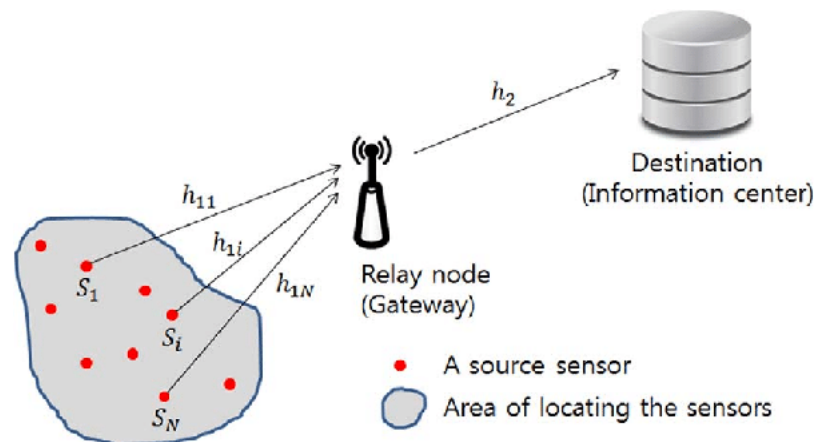


Medical sensor networks



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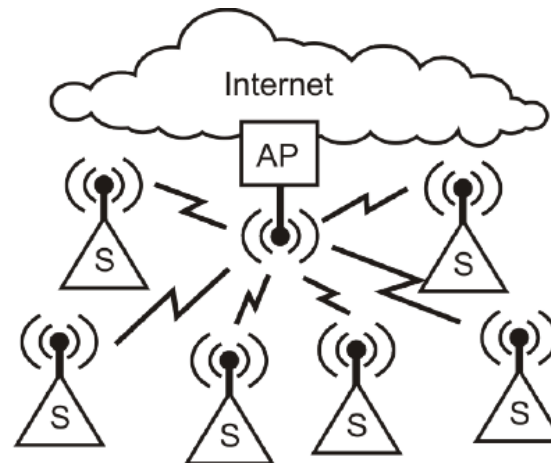
WSN to observe environmental phenomena



Medical sensor networks

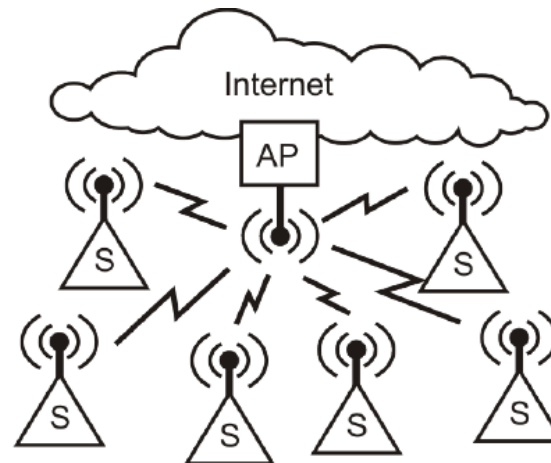
This paper designs an **asynchronized** scheduling policy that achieves the **optimal trade-off** between **Aol** and **Lifetime**

Our System Model



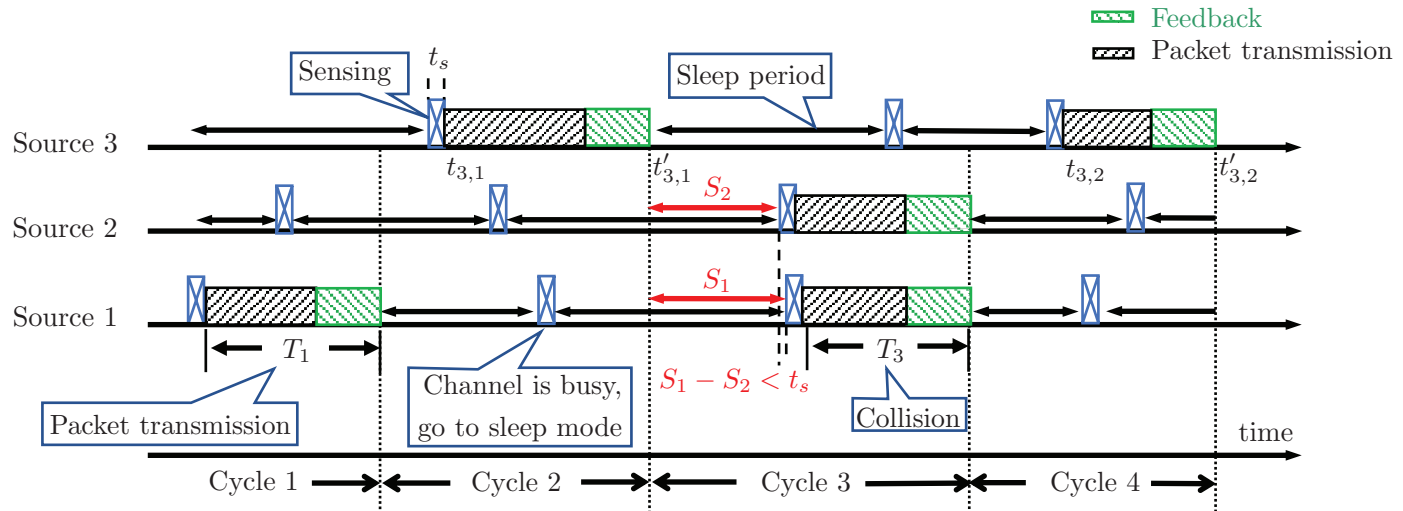
- Wireless network with M sources
- Sources send update packets to an AP via a **shared** channel
- Sources utilize **carrier sensing** to reduce **collisions**

Our System Model



- Wireless network with M sources
- Sources send update packets to an AP via a **shared** channel
- Sources utilize **carrier sensing** to reduce **collisions**
- Each source follows sleep-wake scheme:
 - **Generates** and transmits a **new** packet if the channel is sensed **idle**
 - Sleeps if:
 - Senses the channel to be **busy**
 - Completes a packet transmission

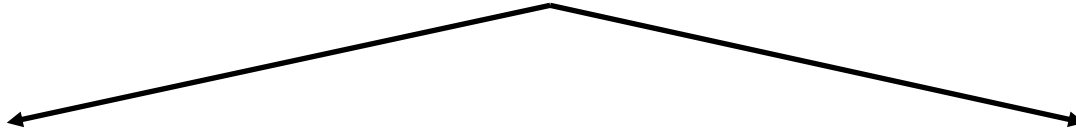
Our System Model (Cont.)



- Source l sleeping period: **Exponentially** distributed with mean $\frac{1}{r_l}$
- Transmission times: **Arbitrarily** distributed with mean $\mathbb{E}[T]$
- Sensing time is t_s
- **Collision** occurs if two sources start transmitting **within a duration** of t_s

Scheduler Types

Scheduling policies

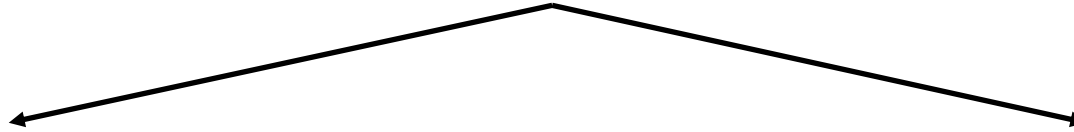


Synchronized schedulers

Asynchronized schedulers

Scheduler Types

Scheduling policies



Synchronized schedulers

Asynchronized schedulers

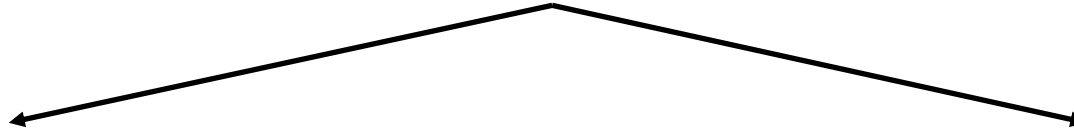
- With access probabilities (e.g., for time-slotted systems)

$$\mathbf{a} = \{a_l\}_{l=1}^M, \sum_{i=1}^M a_i \leq 1$$

- Source l gains channel access after a packet transmission with a probability a_l

Scheduler Types

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Asynchronized schedulers

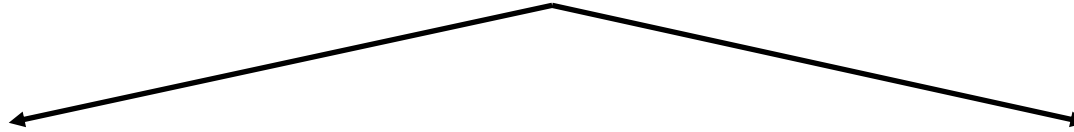
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- **Pros:**
 - Good performance
 - No collision
- **Cons:**
 - Require coordination **overhead**
 - Not implementable in case of:
 - **Dense** networks
 - Non-constant transmission times
- Ex.:[\[Talak, Karaman, Modiano 2018\]](#)

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- Focus of our work
- **Pros:**
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 - Collision occurs due to **non-zero** sensing time
 - Collision increases Aol and energy consumption

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Asynchronized schedulers

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- **Pros:**
 - No coordination overhead
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- **Cons:**
 - Collision occurs due to **non-zero** sensing time
 - Collision increases Aol and energy consumption
- Ex: CSMA to minimize Aol
- [\[Maatouk, Assaad, Ephremides 2019\]](#)
- [\[Wang, Dong 2019\]](#)
 - No energy constraint
 - Zero sensing time
 - Some distributions for transmission times

Q: How to Model Energy Cost of Collisions and Target Lifetime?

- Source l equipped with a Battery with initial level of B_l
- Source l has a **target lifetime** D_l : Minimum time duration before the battery is **depleted**
- Source l Average energy **replenishment** rate R_l
- Maximum allowable energy consumption rate for transmissions

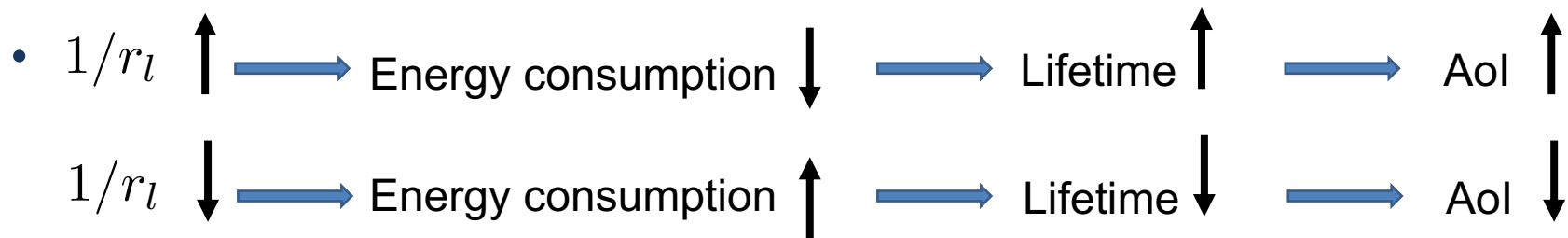
$$E_{\text{con},l} = \frac{B_l}{D_l} + R_l$$

- Source l fraction of time transmitting update packets σ_l
- Source l average energy consumption rate in the transmission mode $E_{\text{avg},l}$

$$\sigma_l E_{\text{avg},l} \leq E_{\text{con},l} \quad \longrightarrow \quad \sigma_l \leq E_{\text{con},l} / E_{\text{avg},l} = b_l$$

b_l : The target energy efficiency of source l

Q: How to Minimize the AoI with Energy (Battery Lifetime) Constraints?



- Target: Design r_l 's s.t.

$$\bar{\Delta}_{\text{opt}}^{\text{w-peak}} \triangleq \min_{r_l > 0} \sum_{l=1}^M w_l \mathbb{E}[\Delta_l^{\text{peak}}]$$

$$\text{s.t. } \sigma_l \leq b_l, \forall l,$$

Δ_l^{Peak} : Peak age of source l

w_l : Weight of source l



$$\bar{\Delta}_{\text{opt}}^{\text{w-peak}} \triangleq \min_{r_l > 0} \sum_{l=1}^M \frac{w_l e^{-r_l \frac{t_s}{\mathbb{E}[T]}}}{r_l} e^{\sum_{i=1}^M r_i \frac{t_s}{\mathbb{E}[T]}} \left(1 + \sum_{i=1}^M r_i \right) + \sum_{l=1}^M w_l$$

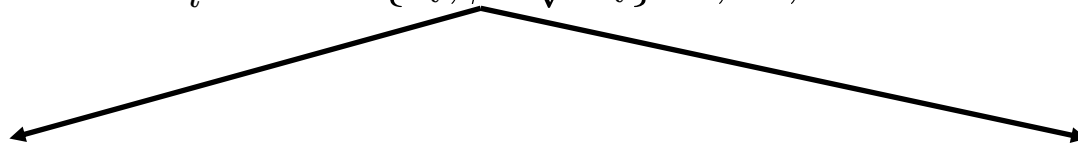
$$\text{s.t. } \frac{[1 - e^{-r_l \frac{t_s}{\mathbb{E}[T]}}] \sum_{i=1}^M r_i + r_l e^{-r_l \frac{t_s}{\mathbb{E}[T]}}}{\sum_{i=1}^M r_i + 1} \leq b_l, \forall l,$$

Non-convex
optimization
problem
(**non-convex**
constraints)

Q: How to Minimize the AoI with Energy (Battery Lifetime) Constraints? (Cont.)

- Derive a low-complexity solution \rightarrow Near-optimal when $\frac{t_s}{\mathbb{E}[T]}$ is small
- $\mathbf{r}^* := (r_1^*, \dots, r_M^*)$

$$r_l^* = \min\{b_l, \beta^* \sqrt{w_l}\} x^*, \forall l,$$



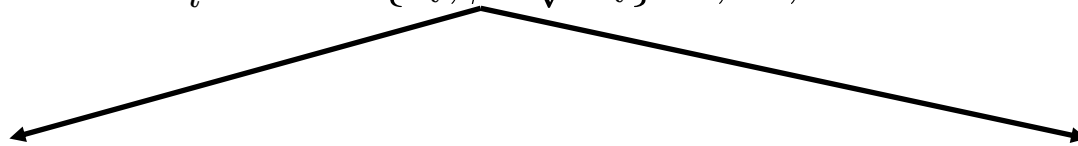
Energy-adequate regime

Energy-scarce regime

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Energy-adequate regime

$$\sum_{i=1}^M b_i \geq 1$$

Energy-scarce regime

- Sufficient energy to ensure that at least one source is awake at any time

β^* : the root of

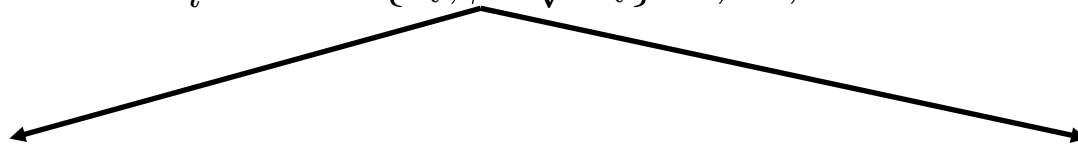
$$\sum_{i=1}^M \min\{b_i, \beta^* \sqrt{w_i}\} = 1$$

$$x^* = \frac{-1}{2} + \sqrt{\frac{1}{4} + \frac{\mathbb{E}[T]}{t_s}},$$

Q: How to Minimize the AoI with Energy (Battery Lifetime) Constraints? (Cont.)

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Energy-scarce regime

$$\sum_{i=1}^M b_i < 1$$

- Sources have to sleep for some time to meet the sources' energy constraints

$$x^* = \frac{\min_l c_l}{1 - \sum_{i=1}^M b_i}, \quad \beta^* = \sum_{i=1}^M \frac{1}{\sqrt{w_i}}$$

$$c_l = \frac{2b_l \left(1 - \sum_{i=1}^M b_i\right)^2}{Q}$$

$$Q = b_l \left(1 - \sum_{i=1}^M b_i\right)^2$$

$$+ \sqrt{b_l^2 \left(1 - \sum_{i=1}^M b_i\right)^4 + 4b_l^2 \left(1 - \sum_{i=1}^M b_i\right)^2 \left(\sum_{i=1}^M b_i - b_l\right) \frac{t_s}{\mathbb{E}[T]}}$$

Main Results

Theorem:

- Our solution is **near-optimal** when $\frac{t_s}{\mathbb{E}[T]}$ is **sufficiently small**, i.e.,

- If $\sum_{i=1}^M b_i \geq 1$, $\left| \bar{\Delta}^{\text{w-peak}}(\mathbf{r}^*) - \bar{\Delta}_{\text{opt}}^{\text{w-peak}} \right| \leq 2\sqrt{\frac{t_s}{\mathbb{E}[T]}} C_1 + o\left(\sqrt{\frac{t_s}{\mathbb{E}[T]}}\right)$

- If $\sum_{i=1}^M b_i < 1$, $\left| \bar{\Delta}^{\text{w-peak}}(\mathbf{r}^*) - \bar{\Delta}_{\text{opt}}^{\text{w-peak}} \right| \leq \frac{t_s}{\mathbb{E}[T]} C_2 + o\left(\frac{t_s}{\mathbb{E}[T]}\right)$

C_1, C_2 : Constants

- Our solution is **asymptotically optimal** as $\frac{t_s}{\mathbb{E}[T]} \rightarrow 0$, i.e.,

$$\lim_{\frac{t_s}{\mathbb{E}[T]} \rightarrow 0} \left| \bar{\Delta}^{\text{w-peak}}(\mathbf{r}^*) - \bar{\Delta}_{\text{opt}}^{\text{w-peak}} \right| = 0$$

Main Results

Corollary: the performance of our proposed algorithm is **asymptotically no worse** than any **synchronized scheduler**, i.e., we have

$$\lim_{\frac{t_s}{\mathbb{E}[T]} \rightarrow 0} \bar{\Delta}_{\text{opt}}^{\text{w-peak}} = \bar{\Delta}_{\text{opt-s}}^{\text{w-peak}}.$$

$\bar{\Delta}_{\text{opt-s}}^{\text{w-peak}}$: Optimal weighted average peak age for synchronized scheduler

Proof Steps

- **Step1:**
 - Check the feasibility of the solution: (Satisfying the energy constraint)
 - Construct an upper bound: Substitute our solution into the obj. function

Proof Steps

- **Step1:**

- Check the feasibility of the solution: (Satisfying the energy constraint)
- Construct an upper bound: Substitute our solution into the obj. function

- **Step 2:**

- Construct a lower bound by relaxing the constraints

$$\min_{r_l > 0} \sum_{l=1}^M \frac{w_l e^{-r_l \frac{t_s}{\mathbb{E}[T]}}}{r_l} e^{\sum_{i=1}^M r_i \frac{t_s}{\mathbb{E}[T]}} \left(1 + \sum_{i=1}^M r_i \right) + \sum_{l=1}^M w_l$$
$$\text{s.t. } \frac{[1 - e^{-r_l \frac{t_s}{\mathbb{E}[T]}}] \sum_{i=1}^M r_i + r_l e^{-r_l \frac{t_s}{\mathbb{E}[T]}}}{\sum_{i=1}^M r_i + 1} \leq b_l, \forall l,$$



$$\min_{r_l > 0} \sum_{l=1}^M \frac{w_l e^{-r_l \frac{t_s}{\mathbb{E}[T]}}}{r_l} e^{\sum_{i=1}^M r_i \frac{t_s}{\mathbb{E}[T]}} \left(1 + \sum_{i=1}^M r_i \right) + \sum_{l=1}^M w_l$$
$$\text{s.t. } r_l \leq b_l \left(\sum_{i=1}^M r_i + 1 \right), \forall l,$$

Proof Steps (Cont.)

- **Step 3:**
 - Analysis the gap between the upper and lower bounds
 - This characterizes the sub-optimality gap of our solution

Easy Implementation

- Uniform solution formula for both energy regimes

$$r_l^* = \min\{b_l, \beta^* \sqrt{w_l}\} x^*, \forall l,$$

- Each source just needs $\beta^* \& x^*$ to compute its sleeping period parameter
- $\beta^* \& x^*$ are functions of $\{(w_i, b_i)_{i=1}^M, t_s/\mathbb{E}[T]\}$

Algorithm 1: Implementation of sleep-wake scheduler.

- 1 The AP gathers the parameters $\{(w_i, b_i)_{i=1}^M, t_s/\mathbb{E}[T]\}$;
 - 2 if $\sum_{i=1}^M b_i \geq 1$ then
 - 3 | The AP derives x^*, β^* according to (19) and (20); \longrightarrow Energy-adequate regime formulas
 - 4 else
 - 5 | The AP derives x^*, β^* according to (25) - (27); \longrightarrow Energy-scarce regime formulas
 - 6 end
 - 7 The AP broadcasts x^*, β^* to all the M sources;
 - 8 Upon hearing x^*, β^* , source l compute r_l^* from (18);
-

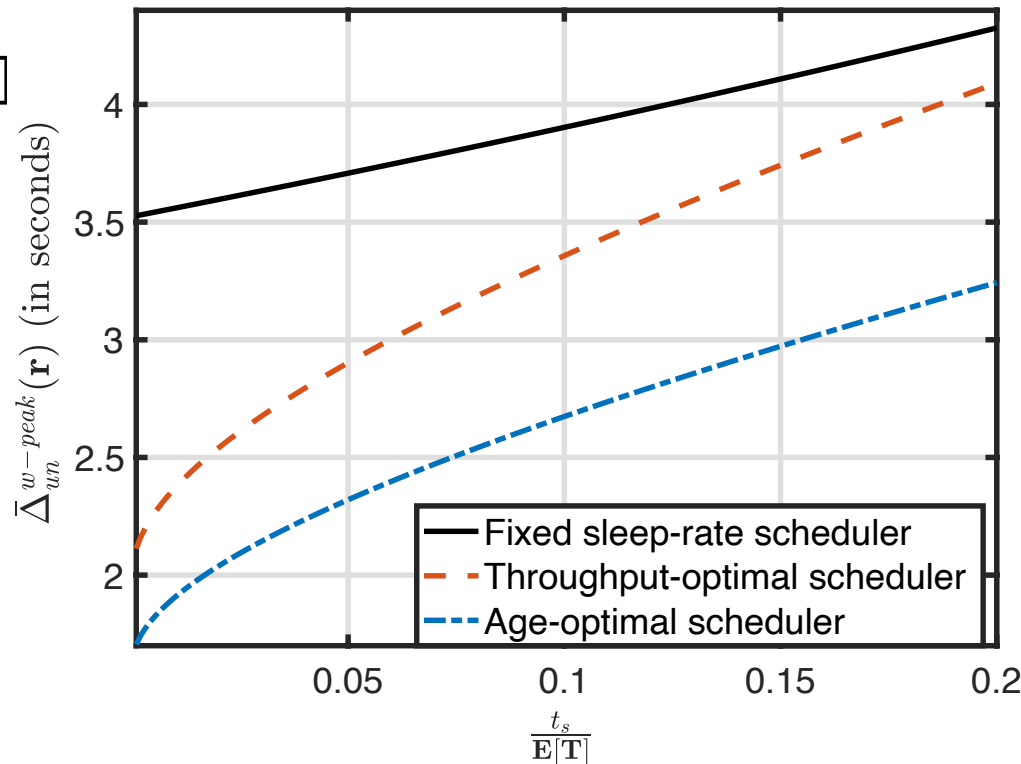
Simulation Results

Fixed-rate: Fixed rates for all sources

Throughput-optimal: Throughput sleep-wake optimal scheduler in [1]

Simulation settings:

- $\mathbb{E}[T] = 5$ ms
- w_l 's are uniformly in $[0, 10]$
- $M = 10$ sources
- b_l 's are uniformly in $[0, 1]$



Observations:

1. t_s increases \rightarrow Packet collision increases \rightarrow AoI increases
2. Age-optimal scheduler outperforms other policies (Throughput optimal scheduler is **not necessarily age-optimal**)

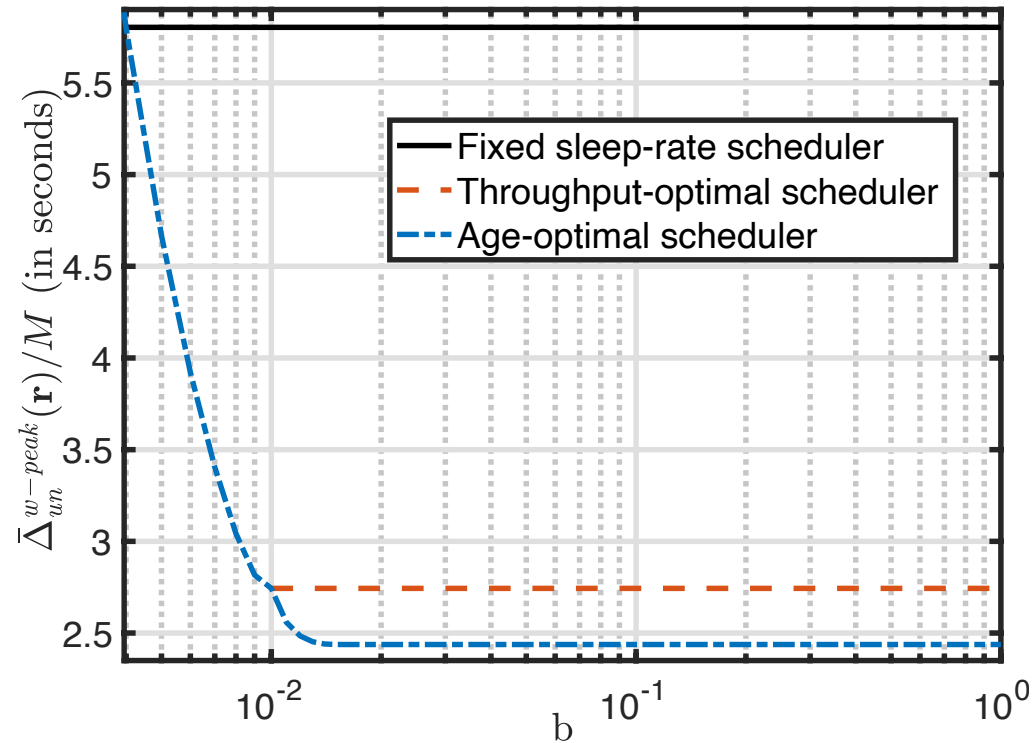
[1] S. Chen, T. Bansal, Y. Sun, P. Sinha, and N. B. Shroff. 2013. Life-Add: Lifetime Adjustable design for WiFi networks with heterogeneous energy supplies. In Proc. WiOpt. 508–515.

Simulation Results (Cont.)

Simulation settings:

- $\mathbb{E}[T] = 5$ ms
- w_l 's are uniformly in $[0, 10]$
- $M = 100$ sources
- $b_l = b \forall l$

Note: Throughput-optimal scheduler is not feasible for $b \leq 0.01$



Observations: $r_l^* = \min\{b_l, \beta^* \sqrt{w_l}\} x^*, \forall l,$

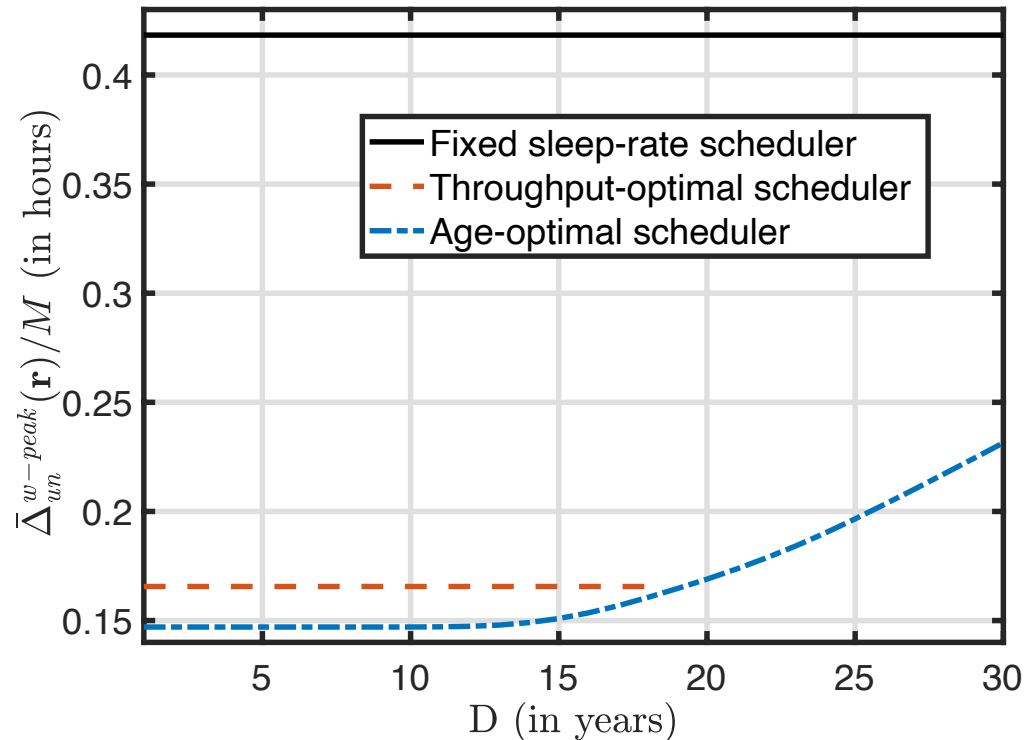
1. There is a value for b after which \mathbf{r}^* is a function solely of w_l 's & β^*
2. Age-optimal scheduler outperforms other policies (Throughput optimal scheduler is **not necessarily age-optimal**)

Simulation Results (Cont.)

Simulation settings:

- $\mathbb{E}[T] = 5$ ms
- w_l 's are uniformly in $[0, 2]$
- $M = 10^5$ sources (**dense** netw)
- $B_l = 8$ mAh, $V_{\text{out}} = 5$ Volt,
 $E_{\text{cons},l} = 24.75$ mW, $\forall l$
- $D_l = D, \forall l$ (Target lifetime)

Note: Throughput-optimal scheduler is not feasible for $D > 18$ years



Observations:

1. D increases \rightarrow Sleeping periods increases \rightarrow Aol increases
2. Age-optimal scheduler can be **active** for **25 years** with a decent average peak age of **0.2 hour**, i.e., **12 minutes**.
3. Age-optimal scheduler outperforms other policies

Summary

- Target: **Efficient** sleep-wake mechanism to attain **optimal trade-off** between minimizing Aol and energy consumption
- Optimization problem is **non-convex**
- Providing a **near-optimal** solution when $\frac{t_s}{\mathbb{E}[T]}$ is a **sufficiently small**
- Providing an easy implementation of our solution
- Our solution is **asymptotically no worse** than any **synchronized scheduler**

Q&A

Thanks

